

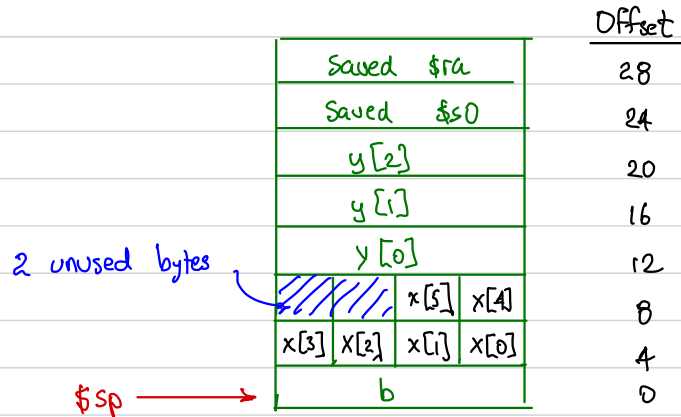
Copied from Monday ...

Stack frame:

C-code:

```
int a, b;
char x[6];
int y[3];

a = g(x, 6);
b = h(y, 3, a);
```



Continuing ... [and midterm #1 next Wednesday, February 12!]

Translations:

```
a = g(x, 6);    addi $a0, $sp, 4    # $a0 = &x[0]
                addi $a1, $zero, 6    # $a1 = 6
                jal g
                add $s0, $zero, $v0    # a = ret. val.
```

```
b = h(y, 3, a); addi $a0, $sp, 12    # $a0 = &y[0]
                addi $a1, $zero, 3    # $a1 = 3
                add $a2, $zero, $s0    # $a2 = a
                jal h
                sw $v0, ($sp)         # b = ret. val.
```

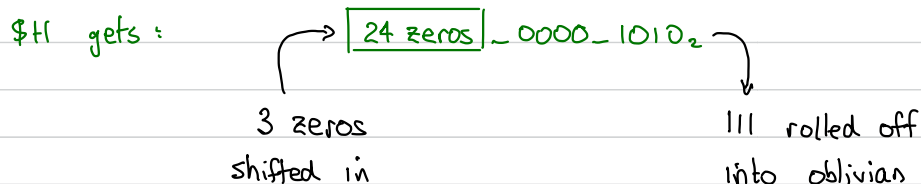
Logical instructions

- sll - shift left logical
- srl - shift right logical

24 zeros

E.g., Let \$t0 = 24 zeros - 0101_0111₂

Suppose: srl \$t1, \$t0, 3



or, ori - logical OR

Recall x OR y

x	y	x OR y
0	0	0
0	1	1
1	0	1
1	1	1

The `or` and `ori` instructions perform 32 simultaneous bit-wise OR operations on 32 pairs of bits:

E.g., Let $\$t0 = 1010_ \boxed{24 \text{ zeros}} _ 0011$
 $\$t1 = 0001_ \boxed{24 \text{ zeros}} _ 0101$

$\$t2 = \underline{1011_ \boxed{24 \text{ zeros}} _ 0111}$

← [Correction, Feb 11!]

`or $t2, $t0, $t1`

ori - one source is a GPR, other is a 16-bit constant

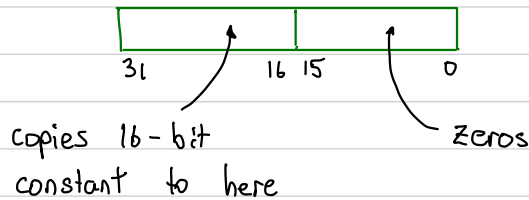
e.g., `ori $t4, $t3, 0x895a`

let $\$t3 = 0xab00_0034$

$\$t3 = 1010_ 1011_ 0000_ 0000_ 0000_ 0000_ 0011_ 0100$
 $\underline{\hspace{10em} 16 \text{ zeros} \hspace{10em}} _ 1000_ 1001_ 0101_ 1010$

$\$t4 = 1010_ 1011_ 0000_ 0000_ 1000 \ 1001_ 0111_ 1110$

lui instruction - load upper immediate



e.g., `lui $t0, 0xf7b3`



$\$t0 = 0xf7b3_0000$

Putting 32-bit constants in a GPR

Can't be done in one instruction - can only do 16 bits at a time.
E.g., a commonly used pseudo-instruction:

```
li    $t1, 0xe9blae09    (load immediate)
```

↑
sticks this 32-bit number into \$t1

Here's what's actually implemented:

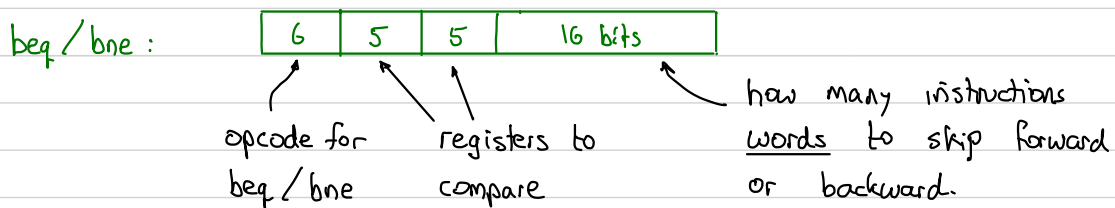
```
lui    $t1, 0xe9bl
ori    $t1, $t1, 0xae09
```

li is identical to la.

and, andi, nor, xor, xori - logical AND, NOR, and XOR functions.

Next: encoding jump and branch instructions

Let's check out the encoding of beq, bne, j, jal, jr.



```
E.g.,    L1:    beq    $t0, $t1, L2    ← (1)
           lw     $t2, ($s0)
           add   $t3, $t3, $t2
           beq   $zero, $zero, L1    ← (2)
L2:      sw     $t3, ($s0)
```

Encoding instruction ①:

$000100 - 01000 - 01001 - 0000 - 0000 - 0000 - 0011_2$
 beq opcode \$t0 \$t1 constant = 3. From the instruction

immediately following the beq (i.e.,
 the lw instruction), count 3 words
 ahead to L2.

$000100 - 01000 - 01001 -$
 beq \$t0 \$t1
 opcode

Encoding instruction ②

$000100 - 00000 - 00000 - 1111 - 1111 - 1111 - 1100_2$
 beq \$zero \$zero

constant = -4 (2's-complement).
 From the sw instruction at L2,
 count 4 words back to L1.