

More ways to compare, <, >

beq, bne only test for equality or inequality, respectively.
How about $\$s0 < \$s1$?

The slt (set on less than) instruction is for this purpose

```

slt  $t0, $s0, $s1      # $t0 = ($s0 < $s1)
bne  $t0, $zero, L1     # if ($t0 != 0) goto L1

```

What the slt instruction does:

```

if ($s0 < $s1)
    $t0 = 1;
else
    $t0 = 0;

```

Using sll (shift-left logical) to multiply by a power of 2

Above, we used

```
sll  $t0, $s3, 2
```

This shifts the contents of $\$s3$ left by 2 bits and places the result in $\$t0$. General form:

```
sll  [dest GPR], [src GPR], [constant]
```

Example: Suppose $\$s0$ contains 24 zeros $_0000_0110_2$ (6_{10})

The instruction: `sll $s1, $s0, 5`

After execution: $\$s0$ unchanged

$\$s1$ 24 zeros $_1100_0000_2$ (192_{10})

→ In decimal, $6 \times 2^5 = 192$

- Notes:
- Bits disappear off the left end
 - zeros shifted in from the right

MIPS pseudoinstructions

pseudo-instructions → line of assembly language text available as a convenience to the programmer

- not real MIPS instructions
- generates one or more real instructions automatically

Example: `blt` (branch on less than)
- produces an `slt` and a `bne` instruction

Warning: In ENCM 369, only real MIPS instructions are usually allowed

Exception: `la` (load address) — as in Lab #2
Very useful!

Procedures in MIPS assembly language

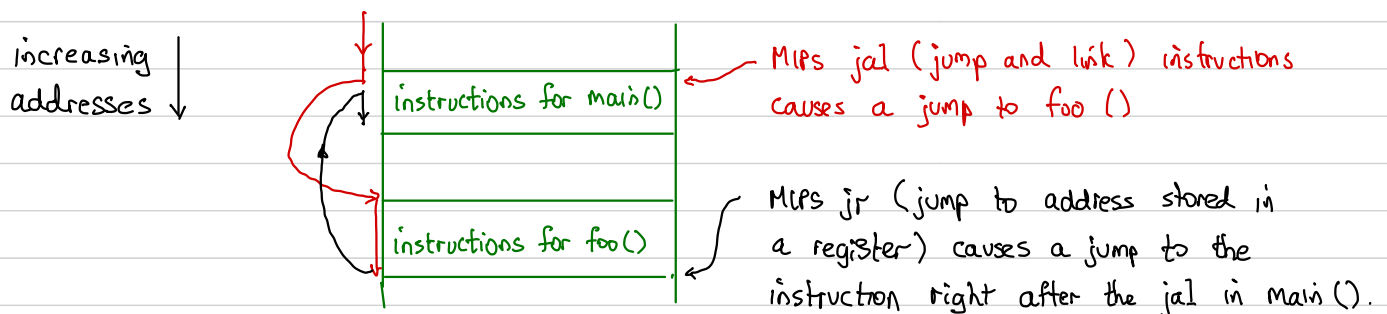
Very important — handles C-language functions

What's involved:

- jumping into and returning from procedures with `jal`, `jr`
- argument passing and return values
- handling register conflicts
- managing the stack to help with
 - register conflicts
 - allocation of local variables

Instruction flow

Suppose we have `main()` and `foo()`, where `main()` calls `foo()`.



The return-address register : \$ra (GPR 31)

From above,

main:

Instructions

jal foo ←
add \$s0, \$v0, \$zero

Call

- put address of add instruction into \$ra
- put address of foo into the PC (program counter)

Instructions

foo:

Instructions

jr \$ra ←

Return

- Copy address in \$ra into the PC.

Summary of jump instructions : j, jal, jr

j label → goto label (no other actions)
jal label → goto label, but remembers address of the next instruction (where to come back) in \$ra.
jr GPR → uses contents of GPR to decide where to jump (like jr \$ra)

Important terms : caller and callee

Above program involves two procedures:

e.g., main (caller)
foo (callee)